

On the terminal costs in LQ games and the application to receding horizon games

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Game theory meets MPC: advances in multi-agent control

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SYSTEMS CONTROL AND MULTIAGENT OPTIMIZATION RESEARCH

1. Introduction
2. From finite to infinite-horizon
3. Implications for receding horizon games
4. Open problems and future directions

- 1. Introduction**
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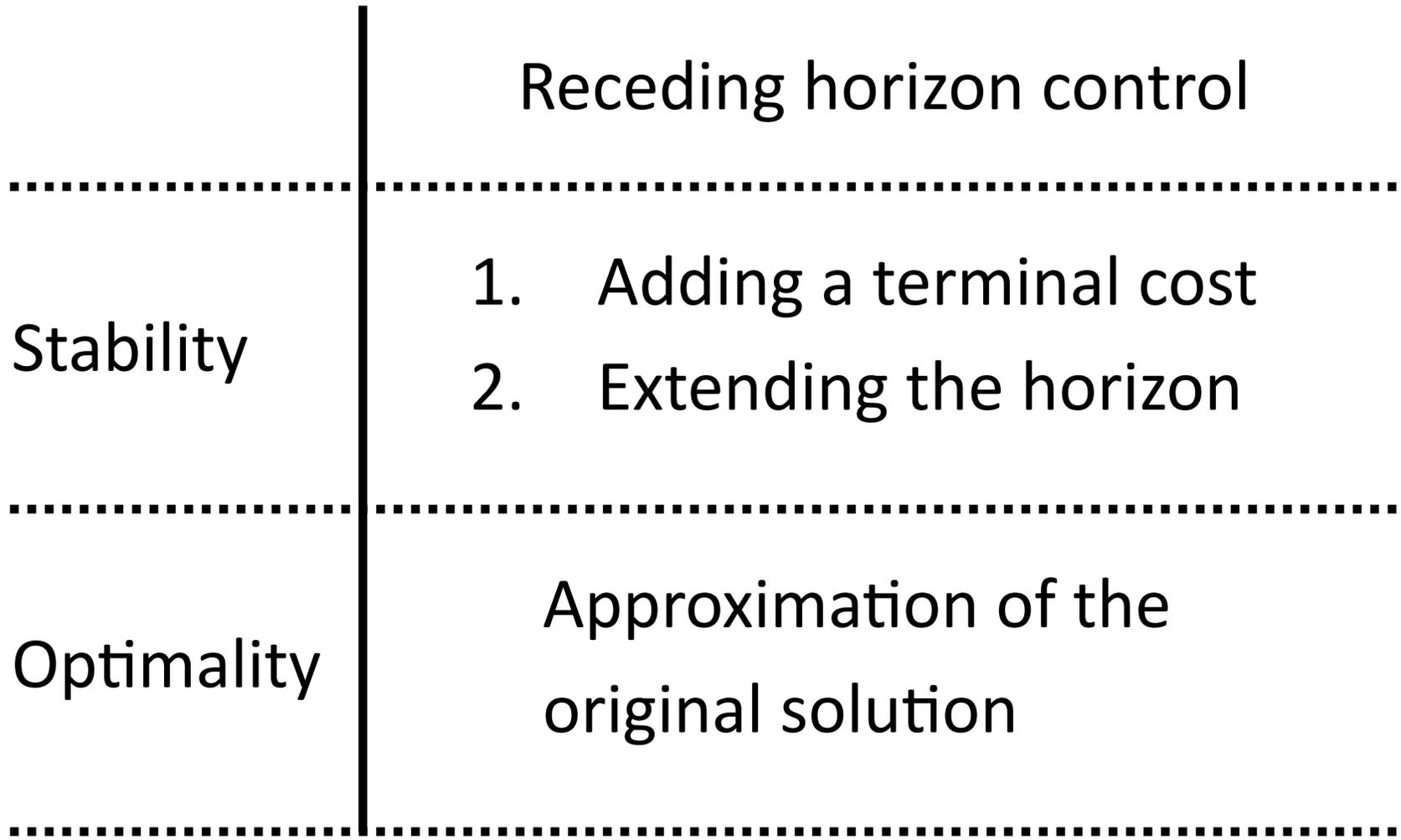
Infinite-horizon control problem

$$\begin{aligned} \min \quad & J(x_0, \mathbf{u}) \\ \text{s.t.} \quad & x_{t+1} = f_t(x_t, u_t) \\ & g_t(x_t, u_t) \end{aligned}$$

The goals are:

- ① Stabilize the system
- ② Minimize the cost function
- ③ Satisfy the constraints

How does receding horizon perform?



Can we extend to receding horizon games?

	Receding horizon control	Receding horizon games
Stability	<ol style="list-style-type: none">1. Adding a terminal cost2. Extending the horizon	
Optimality	Approximation of the original solution	

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Linear quadratic games

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Infinite horizon: there can be multiple Nash equilibria.

Which Nash equilibrium is approximated when using receding horizon? Can we choose?

Main results

1

Extending the horizon is not enough to guarantee stability.

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Terminal costs act as an equilibrium selector.

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Salizzoni, G., Hall, S., & Kamgarpour, M. (2025)
Bridging Finite and Infinite-Horizon Nash Equilibria in Linear Quadratic Games.
arXiv preprint arXiv:2508.20675.

Single agent - Optimal solution for the **finite**-horizon LQR

$$P_t = Q + A^\top P_{t+1} A - A^\top P_{t+1} B (R + B^\top P_{t+1} B)^{-1} B^\top P_{t+1} A, \quad P_T = Q_T$$

$$K_t^* = -(R + B^\top P_t B)^{-1} B^\top P_t A$$



$$P_t = f(P_{t+1})$$

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Multi agent - Nash equilibrium for the **finite**-horizon LQR

$$(P_t^1, P_t^2, \dots, P_t^N) = f(P_{t+1}^1, P_{t+1}^2, \dots, P_{t+1}^N) \quad P_T^i = Q_T^i \quad \forall i \in [1, N]$$

$$(K_t^1, K_t^2, \dots, K_t^N) = g(P_t^1, P_t^2, \dots, P_t^N)$$

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$$(K_t^1, K_t^2, \dots, K_t^N) = g(P_t^1, P_t^2, \dots, P_t^N)$$

$$\mathbf{P}_t = f(\mathbf{P}_{t+1}), \quad \mathbf{P}_T = \mathbf{Q}_T$$

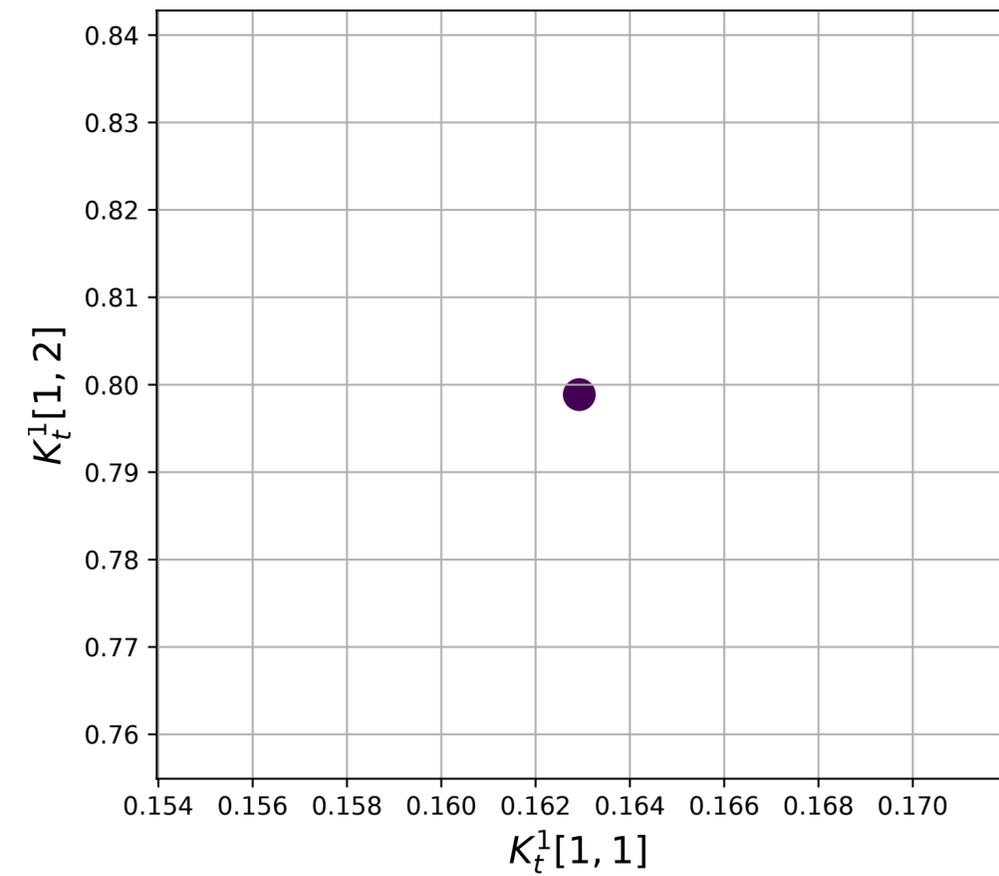
Multi agent - Nash equilibrium for the **infinite**-horizon LQR

$$(P^1, P^2, \dots, P^N) = f(P^1, P^2, \dots, P^N) \quad \longrightarrow \quad \mathbf{P} = f(\mathbf{P})$$

$$(K^1, K^2, \dots, K^N) = g(P^1, P^2, \dots, P^N)$$

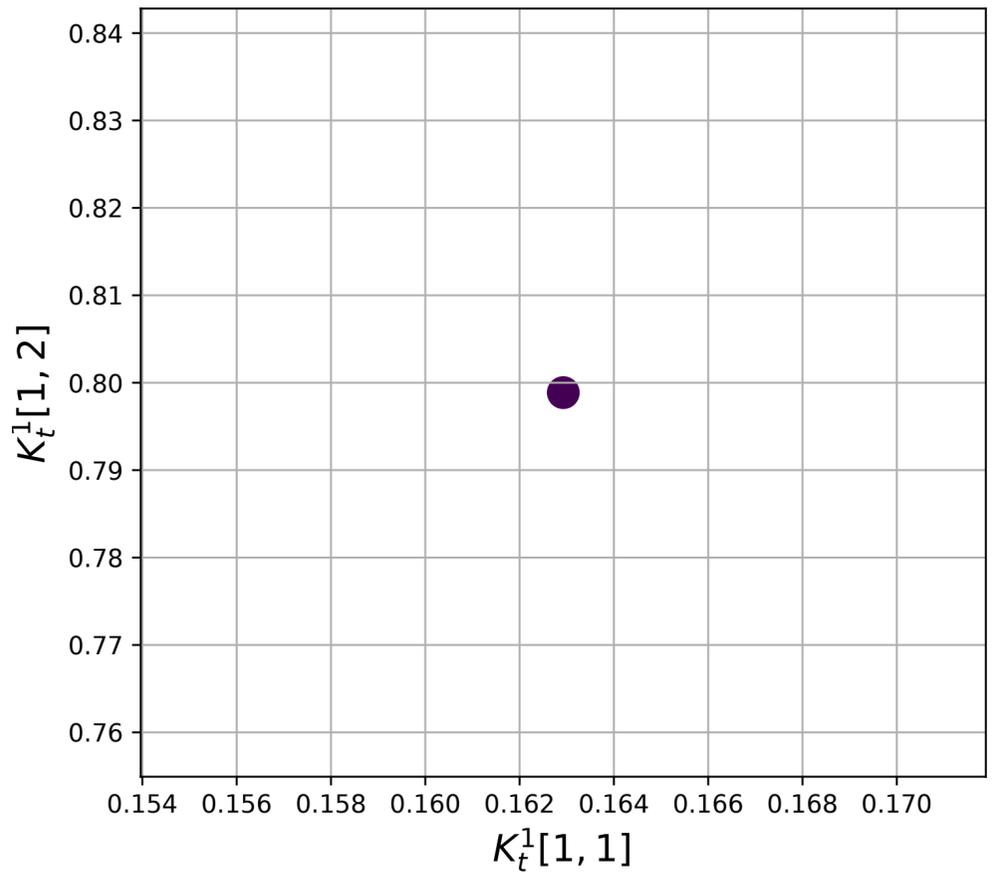
Evolution of the dynamical system $P_t = f(P_{t+1})$

Fixed point

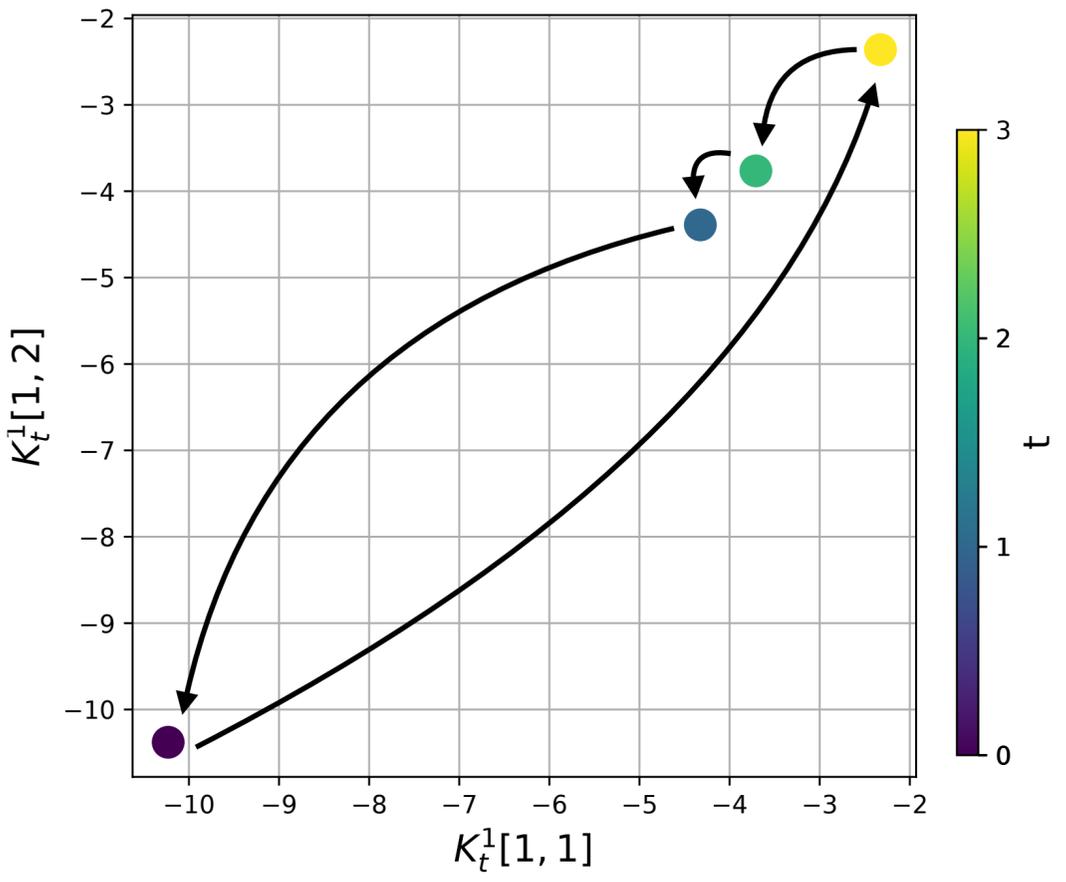


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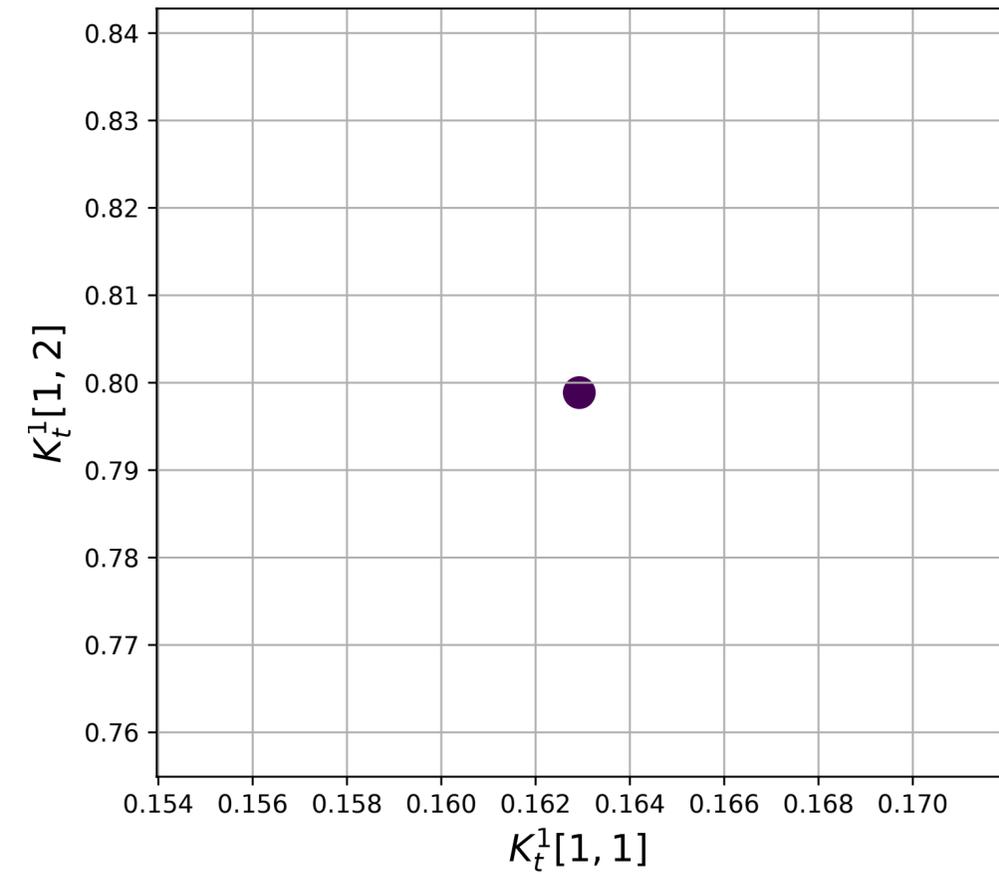


Periodic orbit

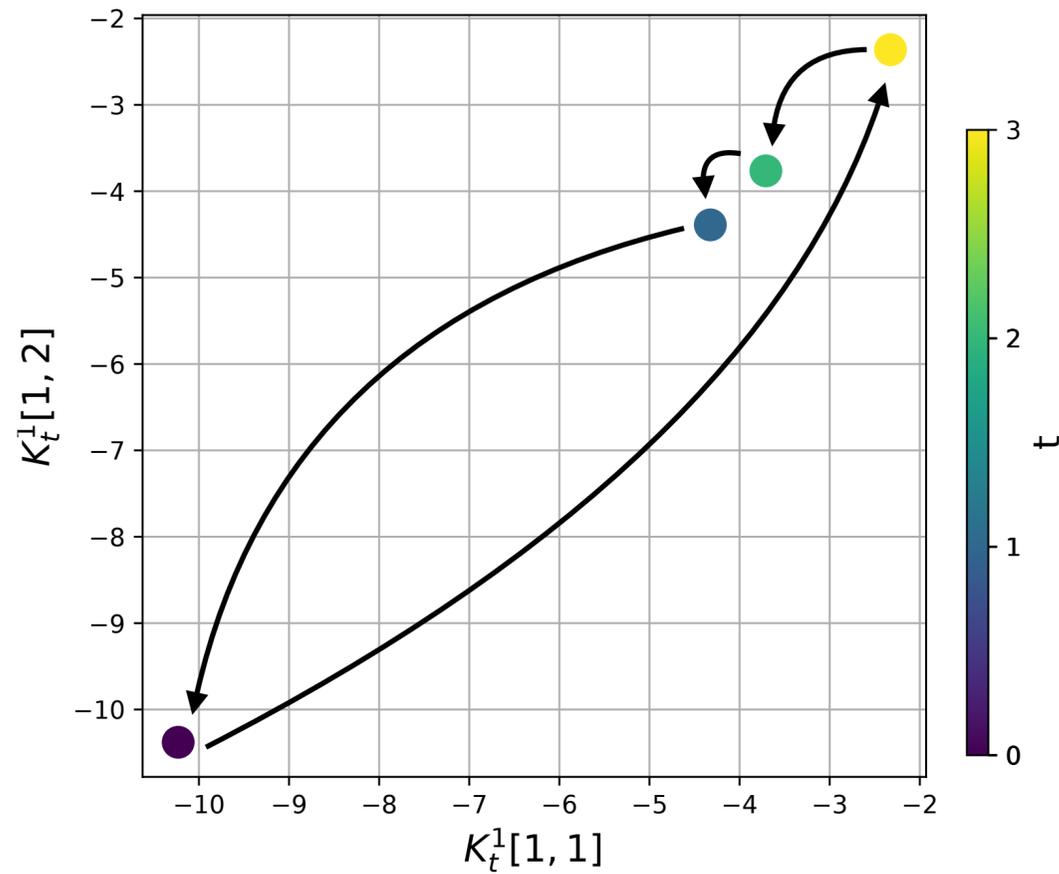


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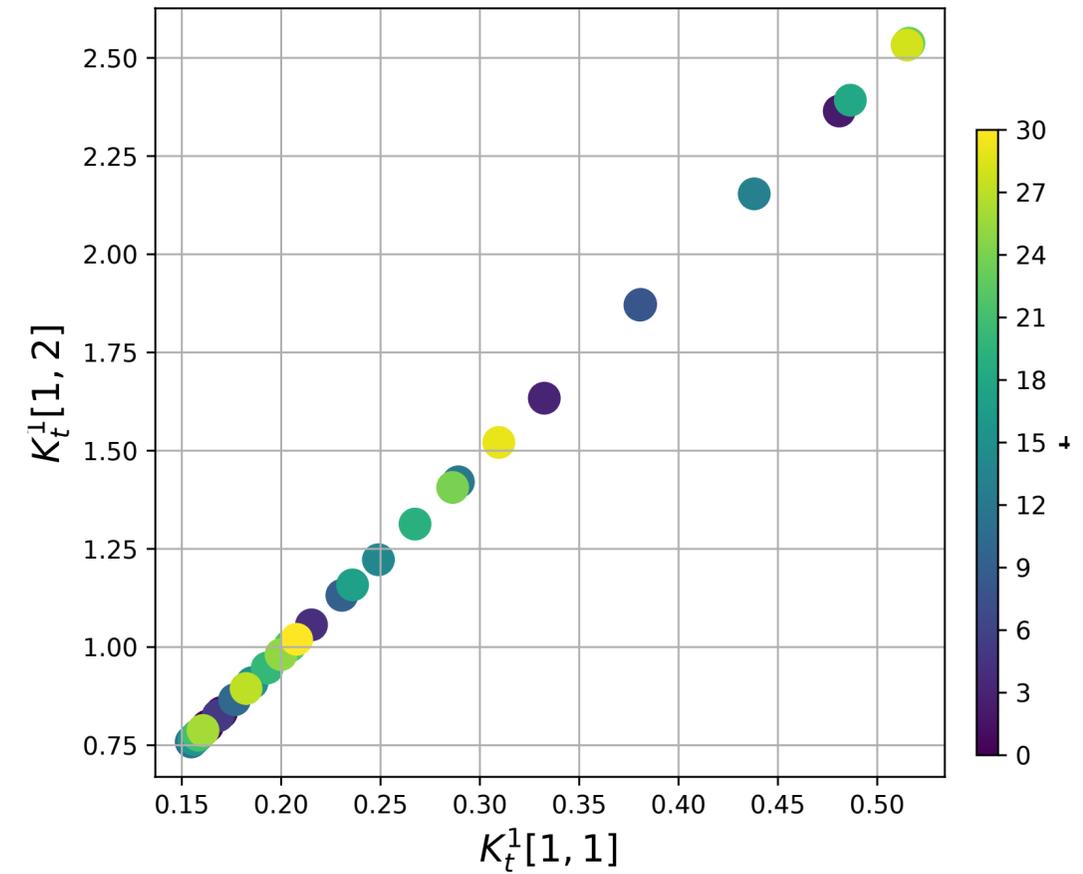
Fixed point



Periodic orbit



Bounded non-convergent



Periodic Nash equilibria

A set of matrices $\{\mathbf{P}_l\}_{l=1}^L$ constitutes a cycle if

$$\mathbf{P}_l = f(\mathbf{P}_{l+1}), \quad l \in [1, L-1], \quad \mathbf{P}_L = f(\mathbf{P}_1).$$

Theorem: Consider a cycle $\{\mathbf{P}_l\}_{l=1}^L$ for a stabilizable game with $Q^i > 0$ for

all $i \in [1, N]$. Then $\{\mathbf{K}_l\}_{l=1}^L$, with $\mathbf{K}_l = g(\mathbf{P}_l)$, constitutes a Nash equilibrium.

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Main results

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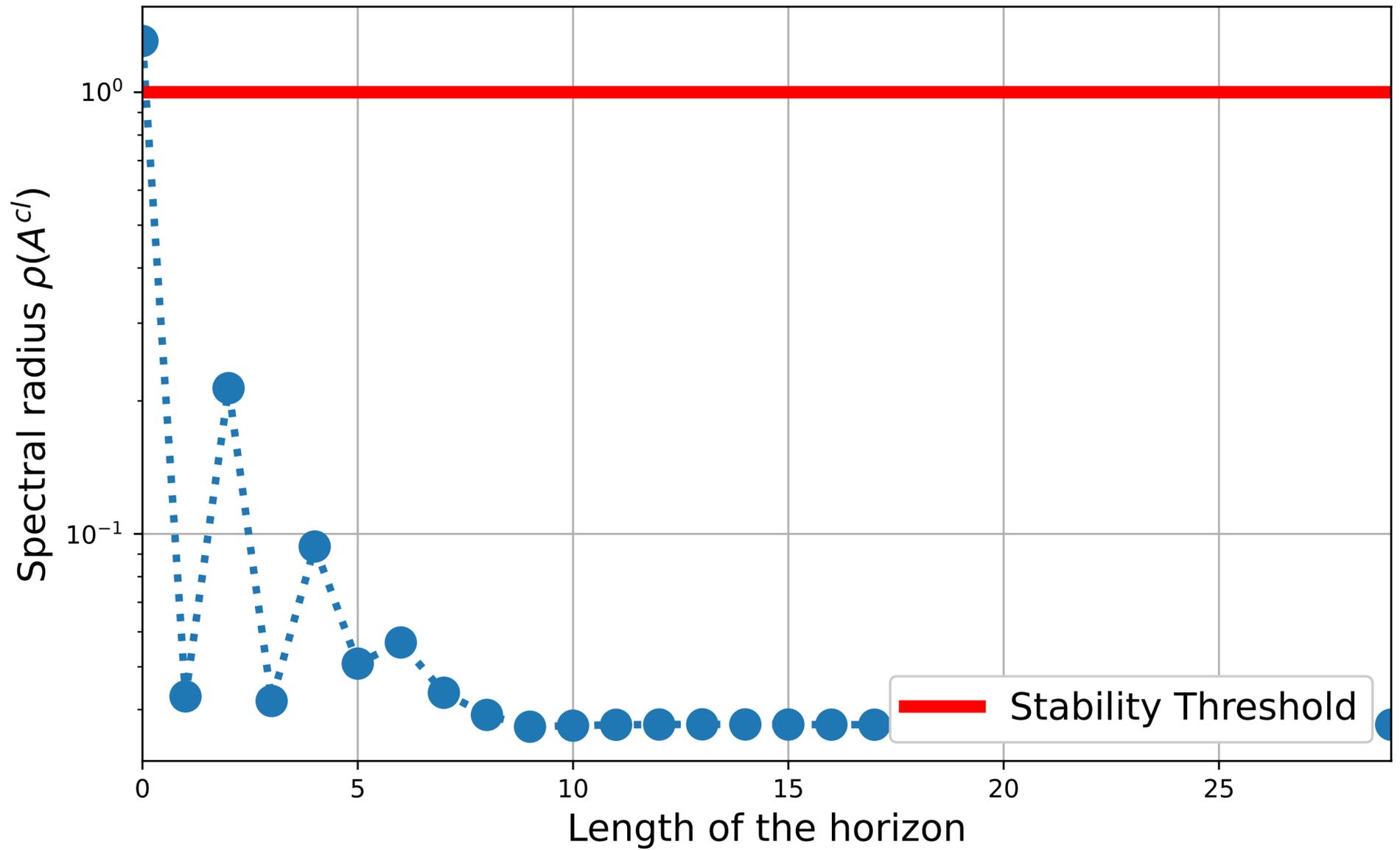
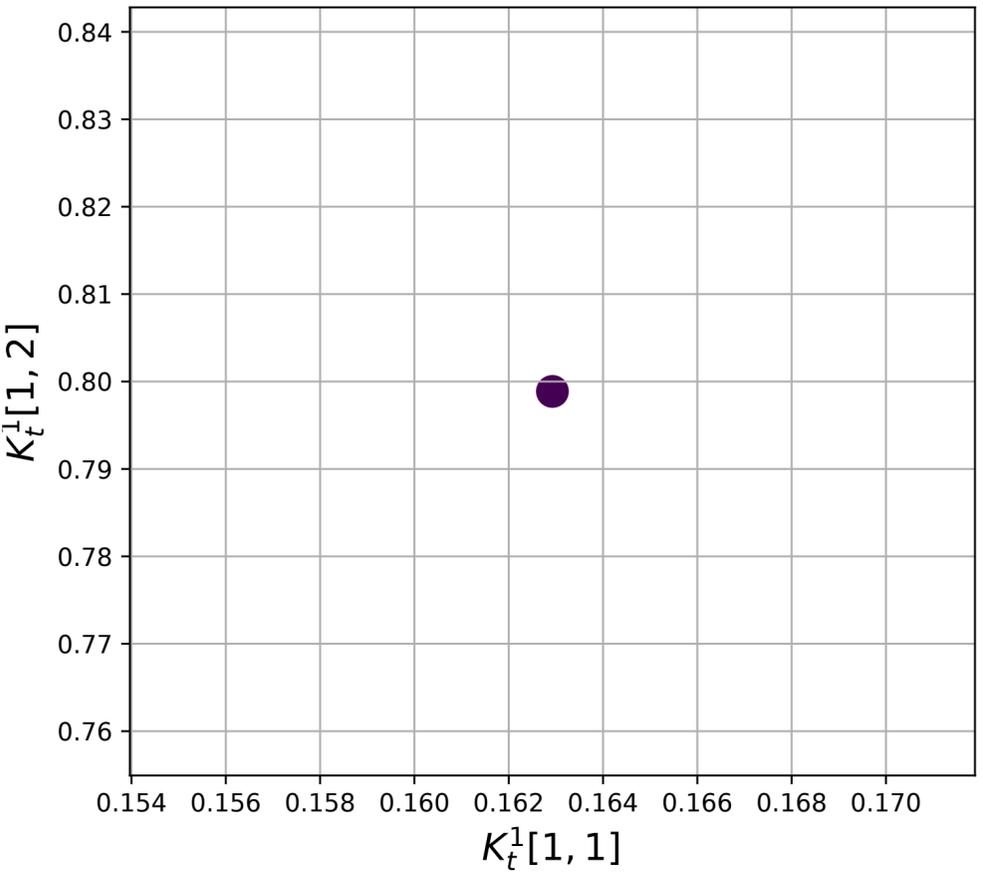
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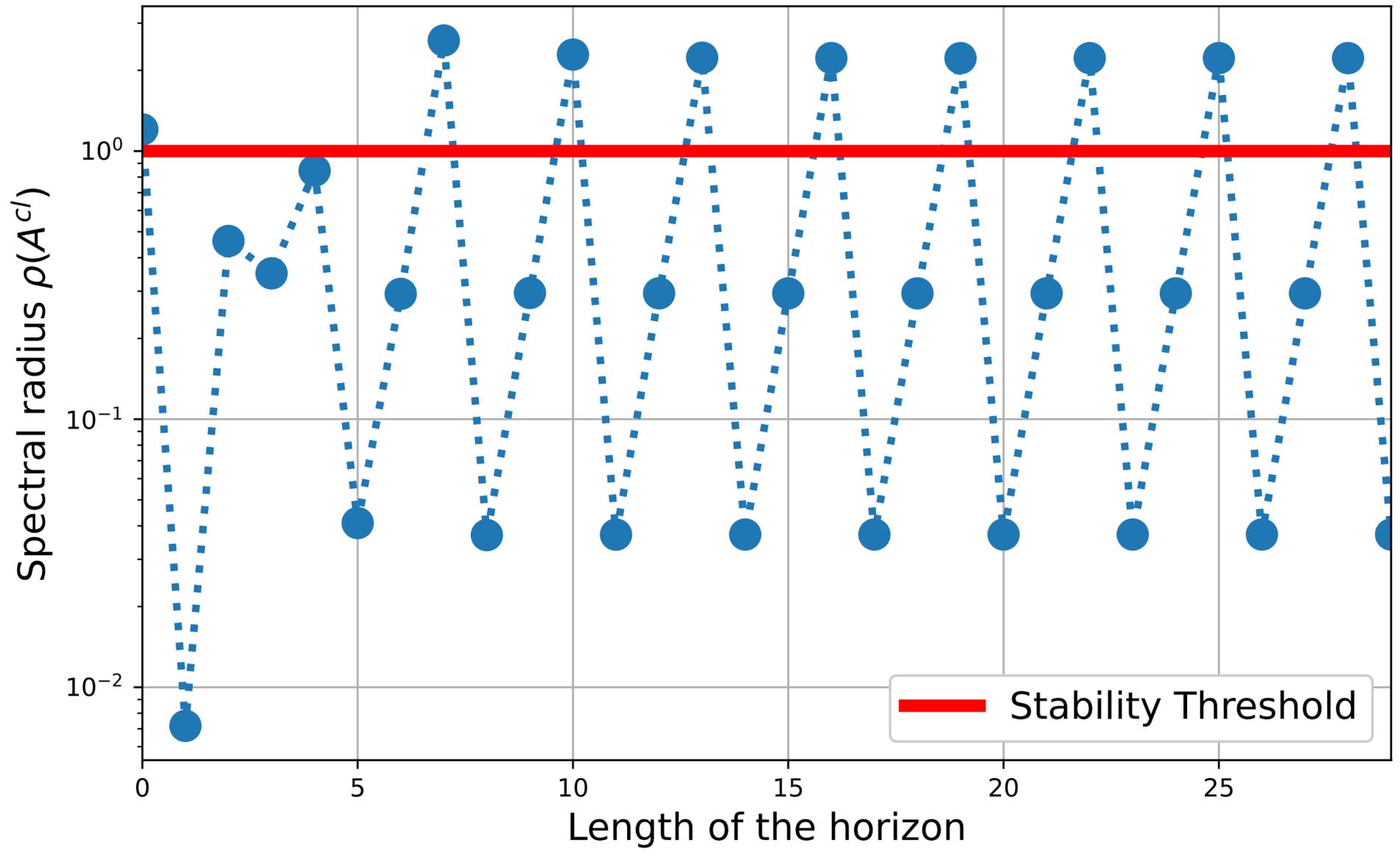
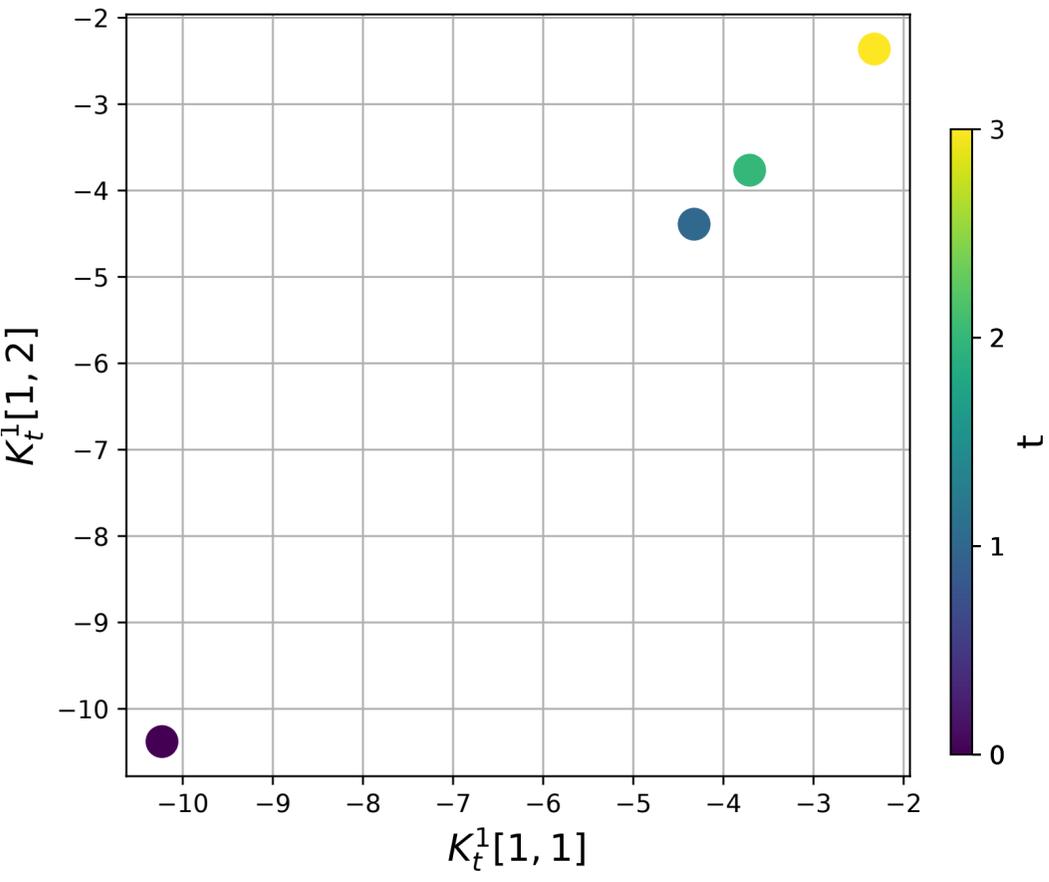
Stability of receding horizon games

Fixed point



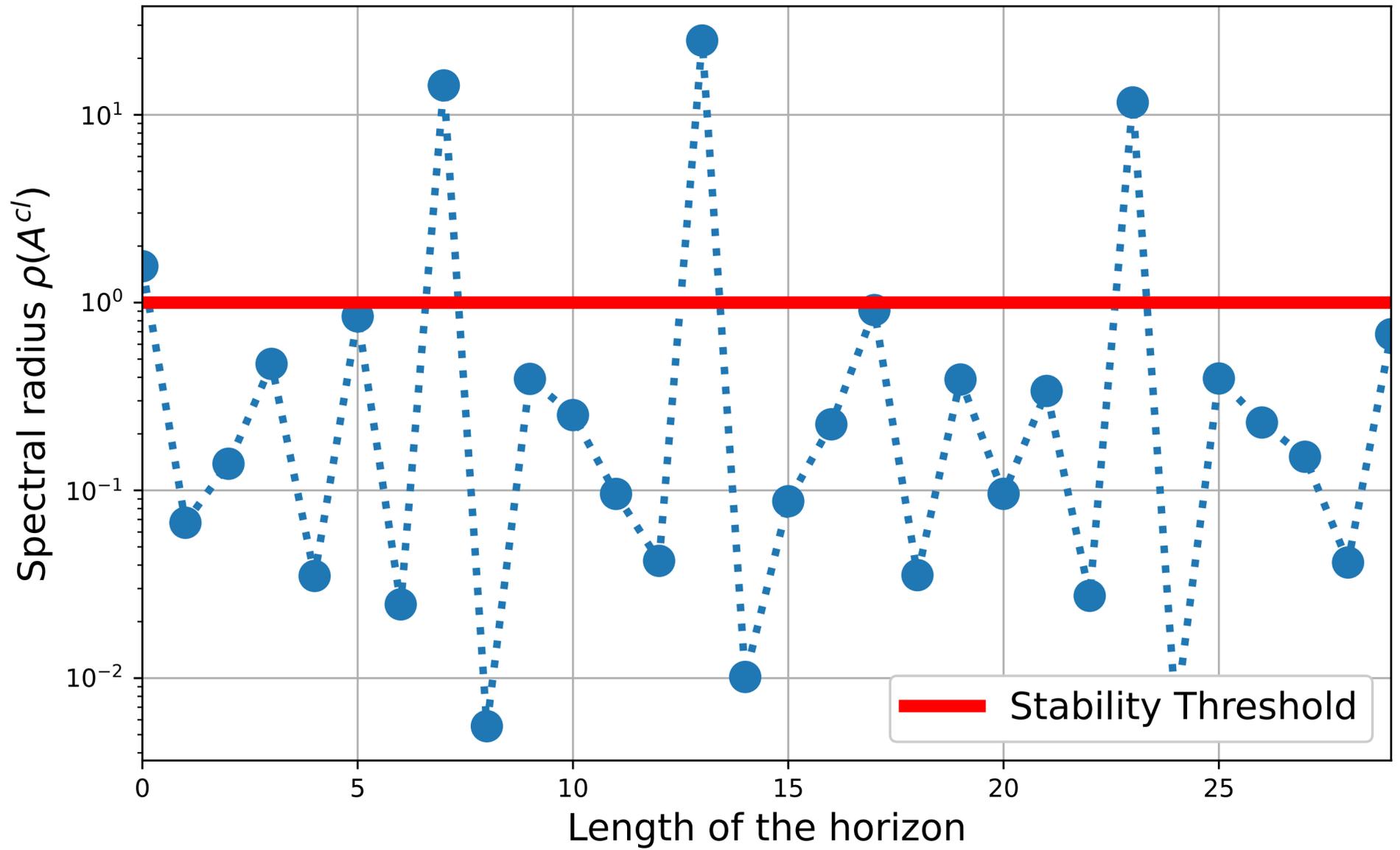
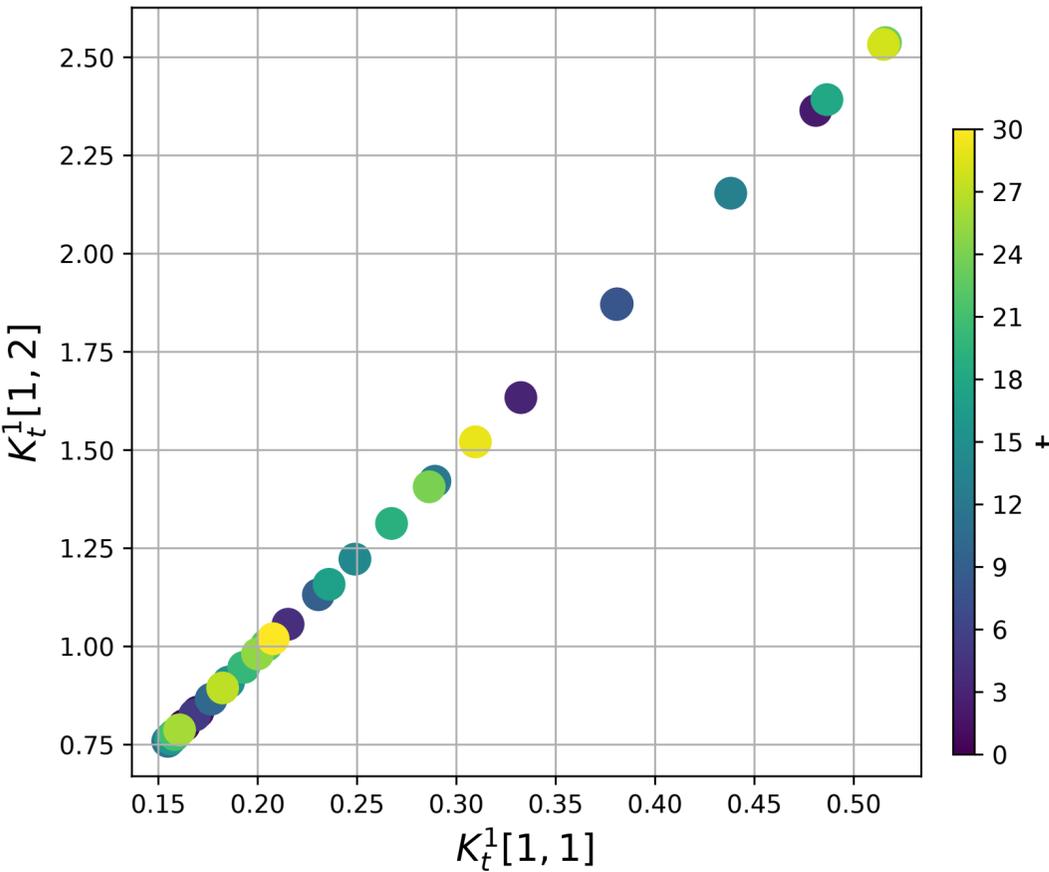
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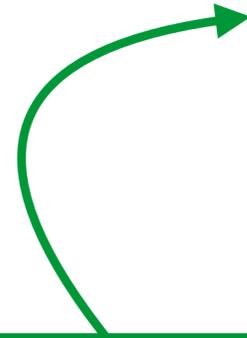


Stability of receding horizon games

Bounded
non-convergent



If the finite-horizon recursion does not converge to a fixed point, the length of the horizon does not matter.



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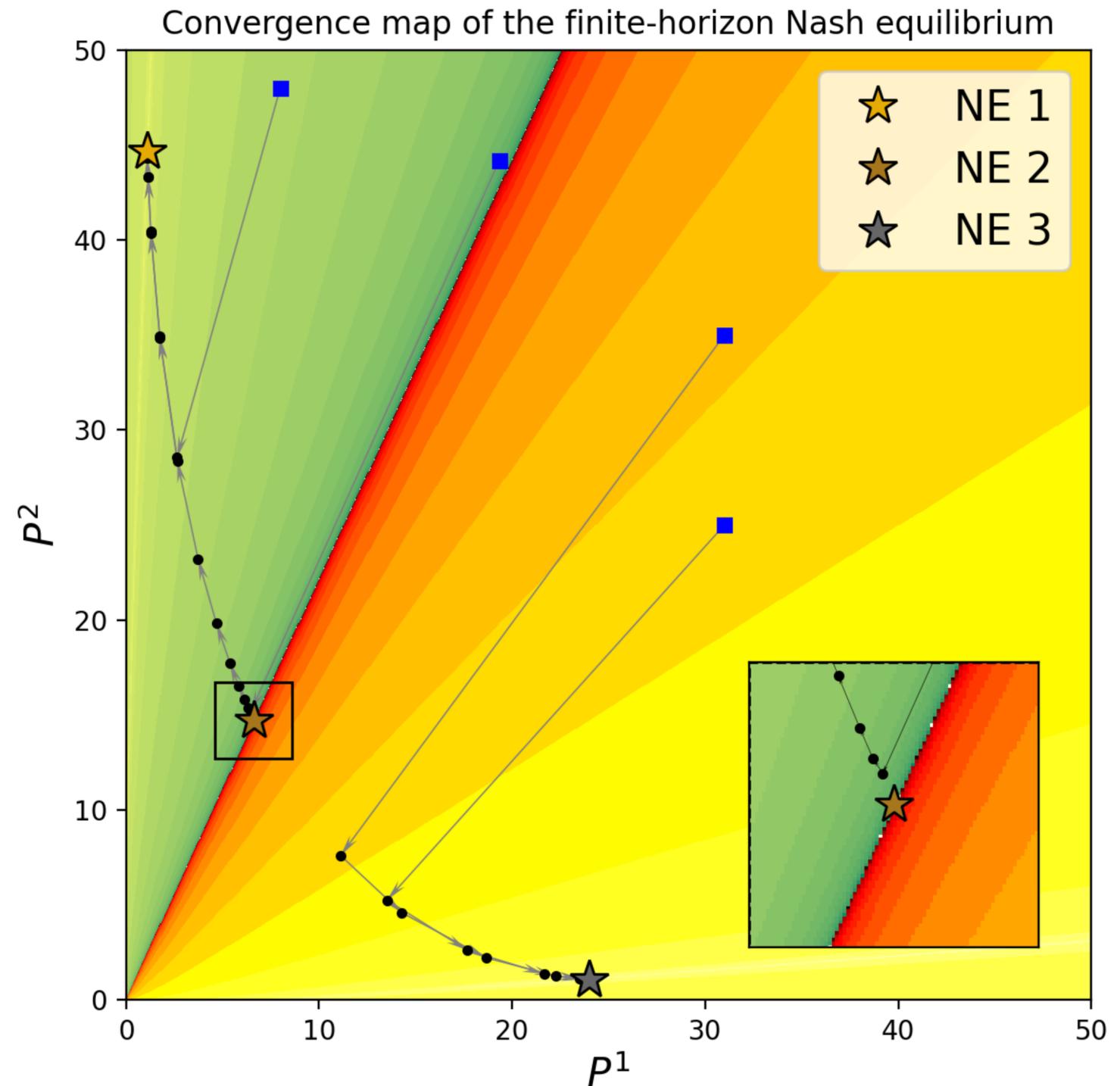
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Attractive and saddle fixed points

Consider a scalar game with two agents and the following setting:

$$A = 3, B_1 = B_2 = 1,$$

$$Q_1 = Q_2 = R_1 = 1, R_2 = 2.$$

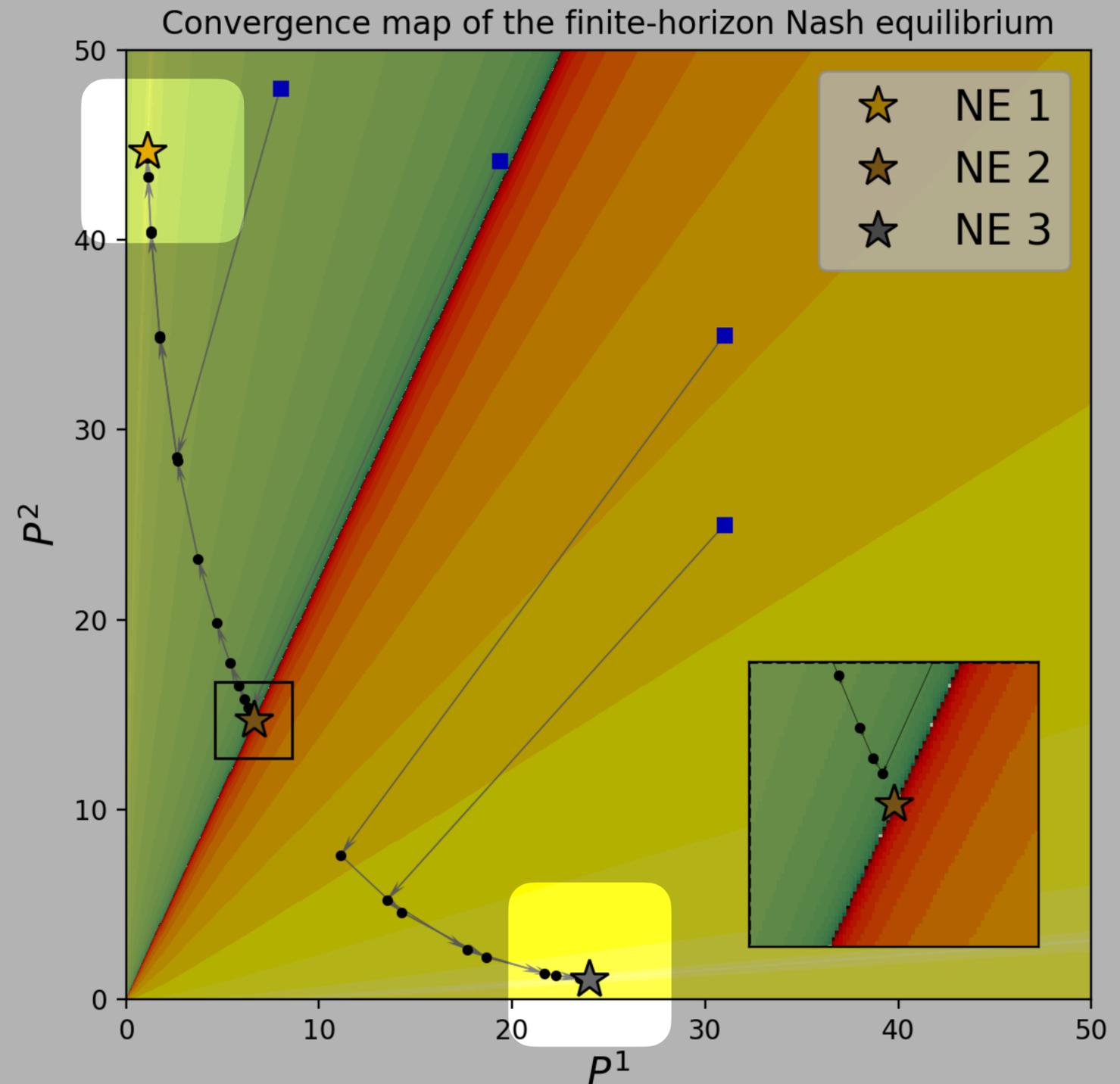


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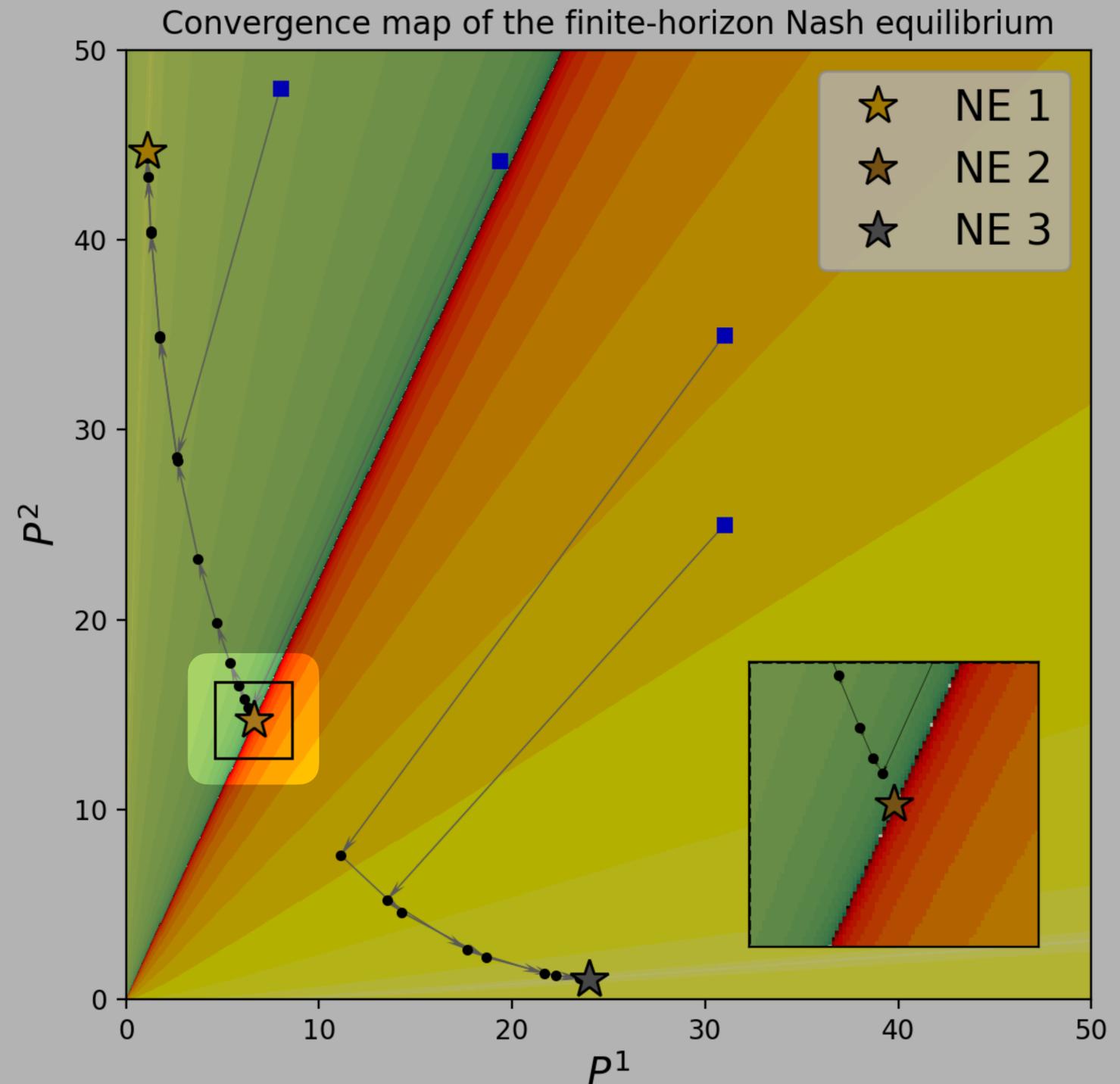


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Not all the Nash equilibria can be approximated. Only the ones attractive for the finite-horizon recursion.

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Open problems and future directions

1. How would constraints affect the results?
2. Can we reformulate the finite-horizon cost to guarantee convergence?
3. How can we compute all the Nash equilibria?